

SubScript: an ACP-based Extension to Scala

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Overview

- Programming is Still Hard
- Algebra of Communicating Processes
- SubScript Now
 - Scala Extension
 - Implementation
 - Debugger demo
- New Dataflow Constructs
 - One-time Flow
 - Lasting Flow
 - SubScript Actors
- Conclusion

Programming is Still Hard

Mainstream programming languages: **imperative**

- good in **batch** processing
- not good in **parsing**, **concurrency**, **event handling**
- Callback Hell

Neglected idioms

- Non-imperative choice: **BNF**, **YACC**
- Data flow: **Unix** pipes
- Process Algebra: **ACP**

Algebra of Communicating Processes - 1

Bergstra & Klop, Amsterdam, 1982 - ...

ACP ~ Boolean Algebra

- + choice
- sequence
- 0 deadlock
- 1 empty process

atomic actions a, b, \dots

parallelism

communication

disruption, interruption

time, space, probabilities

money

...

Algebra of Communicating Processes - 2

Less known than CSP, CCS

Specification & Verification

- Communication Protocols
- Production Plants
- Railways
- Coins and Coffee Machines
- Money and Economy

Strengths

- Familiar syntax
- Precise semantics
- Reasoning by term rewriting
- Events as actions

Algebra of Communicating Processes - 3

$$x+y = y+x$$

$$(x+y)+z = x+(y+z)$$

$$x+x = x$$

$$(x+y) \cdot z = x \cdot z + y \cdot z$$

$$(x \cdot y) \cdot z = x \cdot (y \cdot z)$$

$$0+x = x$$

$$0 \cdot x = 0$$

$$1 \cdot x = x$$

$$x \cdot 1 = x$$

$$(x+1) \cdot y = x \cdot y + 1 \cdot y$$

$$= x \cdot y + y$$

Algebra of Communicating Processes - 4

$$x \parallel y = x \ll y + y \ll x + x | y$$

$$(x+y) \ll z = \dots$$

$$a \cdot x \ll y = \dots$$

$$1 \ll x = \dots$$

$$0 \ll x = \dots$$

$$(x+y) | z = \dots$$

$$\dots = \dots$$

ACP Language Extensions

- 1980: Jan van den Bos - **Input Tool Model**
- 1988-2011: AvD - **Scriptic**
 - Pascal, Modula-2, C, C++, Java
- 1994: Jan Bergstra en Paul Klint - **Toolbus**
- 2011-...: AvD - **SubScript**
 - Scala
 - JavaScript, ... (?)
- Application Areas
 - GUI Controllers
 - Text Parsers
 - Discrete Event Simulation
 - Dataflow Programming (?)
 - Parallel Processing (?)

SubScript Features

"Scripts" – process refinements as class members

```
script a = b; {c}
```

- Much like methods: `override`, `implicit`, named args, varargs, ...
- Invoked from Scala: `_execute(a, aScriptExecutor)`
Default executor: `_execute(a)`
- Body: process expression
Operators: `+` `;` `&` `|` `&&` `||` `/` ...
Operands: script call, code fragment, `if`, `while`, ...
- Output parameters: `?`, ...
- Shared scripts:

```
script send, receive = {}
```

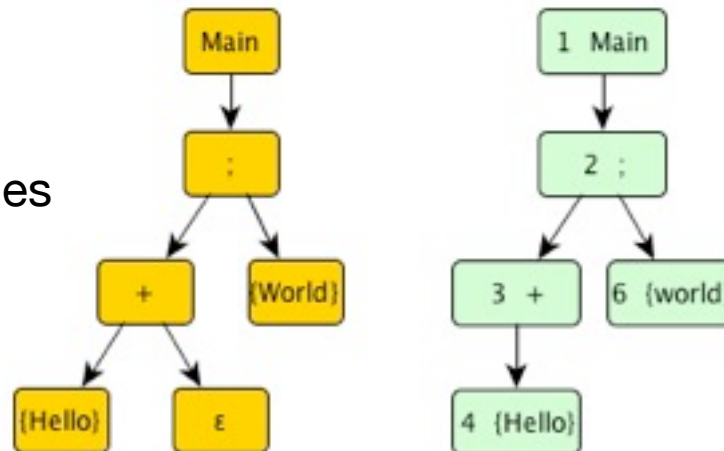
Implementation - 1

- Branch of Scalac: 1300 lines (scanner + parser + typer)

```
script Main = ({Hello} + ε); {World}
```

```
import subscript.DSL._  
def Main = _script('Main) {  
    _seq(_alt(_normal{here=>Hello}, _empty),  
        _normal{here=>World}  
    )  
}
```

- Virtual Machine: 2000 lines
 - static script trees
 - dynamic Call Graph



- Swing event handling scripts: 260 lines
- Graphical Debugger: 550 lines (10 in SubScript)

Debugger - 1

The screenshot displays the Subscript Graphical Debugger interface. At the top, there are control buttons for 'Step', 'Auto', and 'Exit'. The interface is divided into three main sections:

- Call Graph (Top Left):** A tree view showing the current execution context. The root node is 'main', which contains a sub-node ';', which in turn contains two empty object nodes '{}'.
- Log (Bottom Left):** A list of events and messages. The most recent entry is 'Success 4 ;' at line 25. Other entries include 'AAStarted 4 ; 13', 'AAEnded 3 main', and 'Success 5 {}'.
- Control Flow Graph (Right):** A vertical sequence of nodes representing the execution flow: 1 '**', 2 'call', 3 'main', 4 ';', and 5 '{} S'. Node 4 is highlighted with a red box, and a red arrow labeled 'Success' points to it. To the right of node 4, the text 'AA Started' and 'AA Ended' is visible.

Debugger - 2

built using SubScript

```
live = stepping || exit
```

```
stepping = {* awaitMessageBeingHandled *}  
  if (shouldStep)  
    ( @gui: {! updateDisplay !} stepCommand  
      || if (autoCheckBox.isChecked) waitForStep  
        )  
    { messageBeingHandled = false }  
    ...
```

```
exit = exitCommand  
  var isSure = false  
  @gui: { isSure = confirmExit }  
  while (!isSure)
```

```
exitCommand = exitButton + windowClosing
```

One-time Dataflow - 1

```
exit    = exitCommand
        var    isSure = false
        @gui: { isSure = confirmExit }
        while (!isSure)
```

```
exit    = exitCommand; confirmExit ==> while(!_)
```

- Script result type `script confirmExit:Boolean = ...`
- Result values `$ confirmExit^`
`$confirmExit`
- Script Lambda's `<b:Boolean => while(!b)>`
- `x==>y` definition `do_flowTo(<x^>, <y^>)`

```
do_flowTo[T,U](s:script[T], t:T=>script[U]): U = s then t($s)^
```

One-time Dataflow - 2

```
val f: Future[List[String]] = future {session.getRecentPosts}
```

```
f onComplete {  
  case Success(posts) => for (post <- posts) println(post)  
  case Failure(t) => println("An error has occurred: " + t.getMessage)  
}
```

```
f ==> {for (post <- _) println(post)}  
=/> {println("An error has occurred: " + _.getMessage)}
```

Implemented using `implicit script conversion`

- + `==>...=/>`

- + failure exception values (`$$`)

One-time Dataflow - 3

```
val f: Future[List[String]] = future {session.getRecentPosts}
```

```
f onComplete {  
  case Success(posts) => for (post <- posts) println(post)  
  case Failure(t) => println("An error has occurred: " + t.getMessage)  
}
```

```
implicit script future2script[F:Future[F]](f:F)  
= @{f onComplete {case Success(t) => $ = t; there.succeed  
                  case Failure(t) => $$ = t; there.fail   }}: {. .}
```

$x \Rightarrow y \neq z$ definition: `do_flowTo_else(<x^>, <y^>, <z^>)`

```
do_flowTo_else[T,U](s:script[T],  
                    t:T=>script[U],  
                    u:Throwable=>script[U]): U = s then t($s)^ else u($$s)^
```

Lasting Dataflow - 1

```
def copy(in: File, out: File): Unit = {  
  
    val inStream = new FileInputStream(in)  
    val outStream = new FileOutputStream(out)  
  
    val eof = false  
    while (!eof) {  
        val b = inStream.read()  
        if (b==-1) eof=true  
        else outStream.write(b)  
    }  
  
    inStream.close()  
    outStream.close()  
}
```


Lasting Dataflow - 2

```
fileCopier(in:File, out:File) = reader(in) <==> writer(out)
```

```
reader(f:File) = val inStream = new FileInputStream(f);  
                 val b = inStream.read() <=b while (b!=-1);  
                 inStream.close
```

```
writer(f:File) = val outputStream = new FileOutputStream(f);  
                 =>?i: Int while (i!=-1) outputStream.write(i);  
                 outputStream.close
```

```
<==>(i:Int) = {}
```

Lasting Dataflow - 3

```
fileCopier (in:File, out:File) = reader,in &==> writer,out
```

```
fileCrFilter(in:File, out:File) = reader,in &==> crFilter  
                                   &==> writer,out
```

```
crFilter = =>?c:Int if(c!='\r') <=c ...
```

Performance ~ 50 k actions / second

Optimization?

SubScript Actors: Ping Pong

```
class Ping(another: ActorRef) extends Actor {  
  override def receive: PartialFunction[Any,Unit] = {case _ =>}  
    another ! "Hello"  
    another ! "Hello"  
    another ! "Terminal"  
}  
  
class Pong extends SubScriptActor {  
  implicit script str2rec(s:String) = << s >>  
  script ..  
    live = "Hello" ... || "Terminal" ; {println("Over")}  
}
```

SubScript Actors: Partial scripts - 1

```
def receive = {case Request    (r) => sender ! calculate(r)
               case Shutdown   => context.stop(self)
               case Dangerous (r) => a.tell(Work(r), sender)
               case OtherJob   (r) => a!JobRequest(r, sender)
               case JobReply(r,s) => s ! r
            }
```

```
live = .. << case Request    (r) => sender ! calculate(r)
             case Dangerous (r) => a.tell(Work(r), sender)
             case OtherJob   (r) => a!JobRequest(r, sender)
             case JobReply(r,s) => s!r
             >>
; << Shutdown >>
```

SubScript Actors: Partial scripts - 2

```
var initializationReady = false
var activeActors       = 0
var sum: Double        = 0

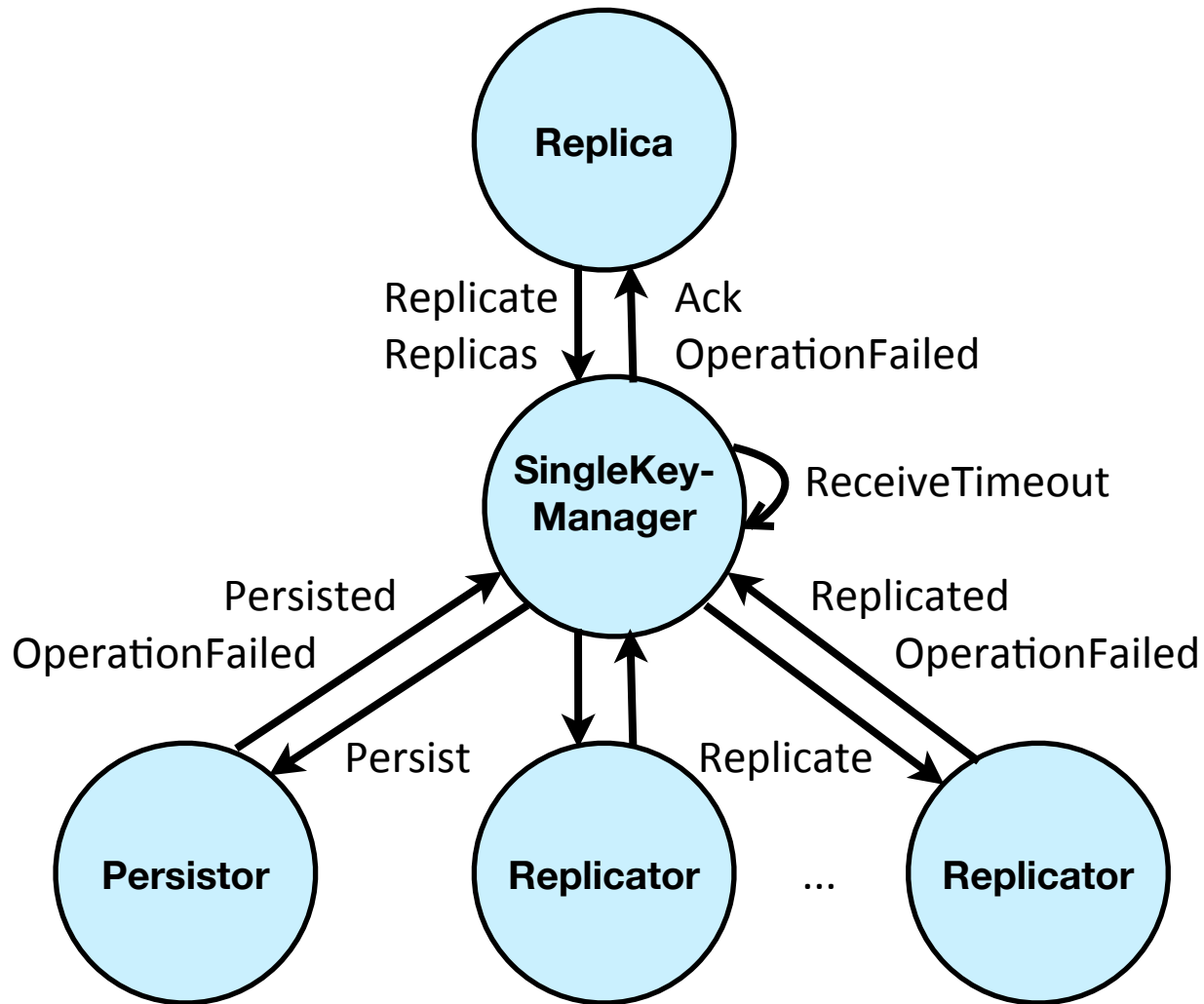
def receive = {
  case context: Context =>
    sum = 0 //reset the instance variables
    activeActors = 0
    for(task <- context.tasks) {
      val actor = actorOf[Delegate].start
      actor ! DoTask(task)
      activeActors += 1
    }
    initializationReady = true

  case delegateResult: Double =>
    sum += delegateResult; sender.get.stop
    activeActors -= 1
    if(initializationReady && activeActors<=0) {
      clientActor ! sum
    }
}
```

SubScript Actors: Partial scripts - 3

```
live = ...
  << context: Context
    ==> var sum: Double = 0
      ( for(task <- context.tasks)
        & {!val actor = actorOf[Delegate].start
          actor ! DoTask(task) !}
        << d:Double => sum+=d; sender.get.stop >>
      )
    {clientActor ! sum}
  >>
```

Akka Actors: KV-Store Assignment - 1



Akka Actors: KV-Store Assignment - 2

```
class SingleKeyManager(key: String) extends Actor { // 44 lines

  var replicatorsOutstandingAck = Set.empty[ActorRef]
  var persistenceOutstandingAck = false
  var outstandingAckId = 0L
  var outstandingValueOption: Option[String] = None
  var ackActor: ActorRef = self // null considered bad
  var ackMsg : AnyRef = self
  var repMsg : AnyRef = self
  var retriesDone = 0

  var timeoutForPersistenceFailure = 100.milliseconds
  var timeoutForPersistenceGiveup = 1.seconds
  var retriesForGiveup = timeoutForPersistenceGiveup / timeoutForPersistenceFailure

  def checkReady() = {
    if (!persistenceOutstandingAck
      && replicatorsOutstandingAck.isEmpty) {
      ackActor ! ackMsg
      context.setReceiveTimeout(Duration.Undefined)
    }
  }
  def sendMsgToPersistorAndReplicators() = {
    if (persistenceOutstandingAck)
      persistor ! Persist(key, outstandingValueOption, outstandingAckId)
    replicatorsOutstandingAck.foreach(_ ! repMsg)
  }

  def receive: Receive = ~~~
}
```


Akka Actors: KV-Store Assignment - 3

```
class SingleKeyManager(key: String) extends Actor { // 44 lines
  ~~~
  def receive: Receive = LoggingReceive{
    case (r@Replicate(key, valueOption, id), ackActor: ActorRef, ackMsg: AnyRef) =>
      context.setReceiveTimeout(timeoutForPersistenceFailure)
      outstandingAckId = id
      outstandingValueOption = valueOption
      persistenceOutstandingAck = true
      replicatorsOutstandingAck = replicasToReplicatorsMap.values.toSet
      retriesDone = 0
      this.ackActor = ackActor
      this.ackMsg = ackMsg
      this.repMsg = r
      sendMsgToPersistorAndReplicators()

    case Persisted (key, msgId) => if (msgId==outstandingAckId)
      {persistenceOutstandingAck = false; checkReady()}
    case Replicated(key, msgId) => if (msgId==outstandingAckId)
      {replicatorsOutstandingAck -= sender; checkReady()}
    case Replicas (replicas) => replicatorsOutstandingAck &=
      replicasToReplicatorsMap.values.toSet; checkReady()

    case ReceiveTimeout =>
      if (retriesDone < retriesForGiveup - 1) {
        retriesDone += 1
        sendMsgToPersistorAndReplicators()
      }
      else ackActor ! OperationFailed(outstandingAckId)
  }
}
```

Akka Actors: KV-Store Assignment - 4

```
class SingleKeyManager2(key: String) extends SubScriptActor { // 44 lines

  var persistorTBD           = true
  var replicatorsTBAck       = Set.empty[ActorRef]
  var valueOption: Option[String] = None
  var tbAckId                = -1L

  val timeoutForPersistenceFailure = 100.milliseconds
  val timeoutForPersistenceGiveup  = 1.seconds
  val retriesForGiveup             = timeoutForPersistenceGiveup / timeoutForPersistenceFailure

  def keyIsStored = valueOption != None
  def replicas    = replicasToReplicatorsMap.values.toSet // from parent class

  def script..

    live = handleReplicateMsg / ..
          || handleReplicasMsg ...
          ; while(keyIsStored)

    handleReplicateMsg = ~~~

    handleReplicasMsg  = ~~~
}
```

Akka Actors: KV-Store Assignment - 5

```
class SingleKeyManager2(key: String) extends SubScriptActor { // 44 lines
  ~~~
  def script..
    live = ~~~
    handleReplicateMsg = << Replicate (r@Replicate(key, valueOption, id),
      replyActor: ActorRef, ackMsg: AnyRef)
      => this.tbAckId      = id
         this.valueOption  = valueOption
         this.replicasTBAck = replicas
         this.persistorTBD = true
         var replyMsg      = OperationFailed(tbAckId)

      ==> ( times(retriesForGiveup);
           timeout(timeoutForPersistenceFailure)
           || ( delegatePersistor & delegateReplicators(r) )
              {!replyMsg=ackMsg!}
              break_up2
           )
         ; {!replyActor!replyMsg!}
      >>

    handleReplicasMsg = ~~~
    delegatePersistor = ~~~
    delegateReplicators(replicateMsg: Any) = ~~~
}
```

Akka Actors: KV-Store Assignment - 6

```
class SingleKeyManager2(key: String) extends SubScriptActor { // 44 lines
  ~~~
  def script..

  live = ~~~

  handleReplicateMsg = ~~~

  handleReplicasMsg = <<Replicas(rs) => foreach (r< replicasTBack-rs)
                        {self!Replicated(key,tbAckId)} >>

  delegatePersistor = if (persistorTBD) (
                        {persistor ! Persist(key, valueOption, tbAckId)})
                        << Persisted (key, @tbAckId) => persistorTBD=false >>
                    )

  delegateReplicators(replicateMsg: Any)
    = for (r <- replicasTBack)
      & {r ! replicateMsg}
      << Replicated(key, @tbAckId) if sender==r
        => replicasTBack -= r >>
}
}
```

Conclusion

- Easy and efficient programming
- Support in Scalac branch
- Simple implementation: 5000 lines, 50%
- Still much to do and to discover
- Open Source:
subscript-lang.org
github.com/AndreVanDelft/scala
- **Help is welcome**
Participate!

The End

- Spare Slides next

SubScript Features - 1

"Scripts" – process refinements as class members

- Called like methods from Scala
 - with a `ScriptExecutor` as extra parameter
- Call other scripts
- Parameters: in, out?, constrained, forcing

Formal	<code>implicit key(c??: Char) = ...</code>		
Actual Calls	Output	Constrained	Forcing
Conventional	<code>key(c?)</code>	<code>key(c? if? c.isDigit)</code>	<code>key('1')</code>
No parentheses	<code>key,c?</code>	<code>key,c? if? c.isDigit</code>	<code>key,'1'</code>
Using <code>implicit</code>	<code>c?</code>	<code>c? if? c.isDigit</code>	<code>'1'</code>

SubScript Features - 2

ACP Atomic Actions ~ Scala Code {...} start/end

{ ... }	Normal
{? ... ?}	Unsure
{! ... !}	Immediate
{* ... *}	New thread
@gui: { ... }	GUI thread
@dbThread: { ... }	DB thread
@reactor: {.}	Event handler
@reactor: {... ... }	Event handler, permanent
@startTime: { ... }	Simulation time + real time
@processor=2: {* ... *}	Processor assignment
@priority=2: {* ... *}	Priority
@chance=0.5: { ... }	Probability

SubScript Features - 3

N-ary operator	Meaning
; WS	Sequence
+	Choice
&	Normal parallel
	Or-parallel
&&	And-parallel
	Or-parallel
==>	Network/pipe
/	Disrupt
%/	Interrupt

Unary operator	Meaning
x*	Process launch

Construct	Meaning
here	Current position
@ ... :	Annotation
if-else	
match	
try-catch-	
for	
while	
break	
...	while(true)
..	Both ... and .
.	Optional break
(-), (+),	Neutral: 0, 1-like

SubScript Features - 4

Process Communication

Definitions: Shared Scripts

`send(i:Int), receive(j??:Int) = {j=i}`

`send(i:Int), receive(i??: _) = {}`

`ch<-(i:Int), ch->(i??:Int) = {}`

`ch<-->(i??:Int) = {}`

`<-->(i??:Int) = {}`

`<==>(i??:Int) = {}`

Usage: Multicalls

<code>send(10) & receive(i?)</code>	Output param
---	--------------

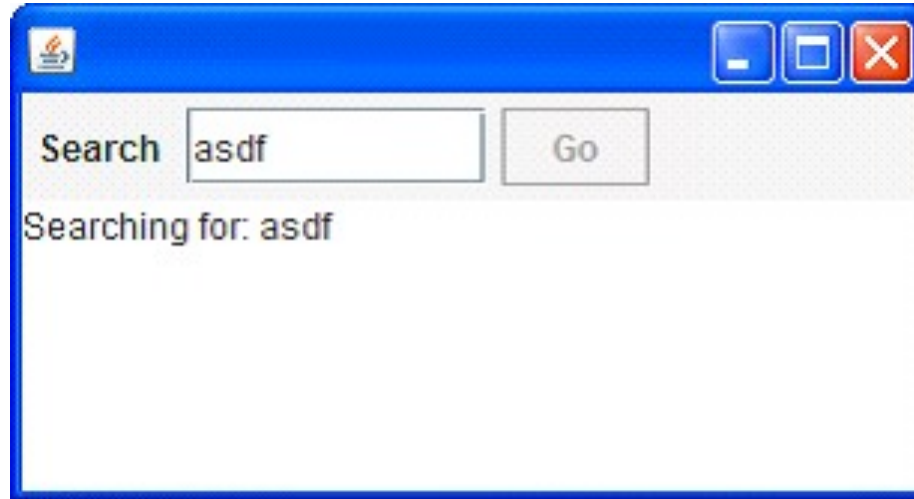
<code>send(10) & receive(10)</code>	Forcing
---	---------

<code>ch<-(10) & ch->(10)</code>	Channel
--	---------

<code><-10 & ->i?</code>	Nameless
------------------------------------	----------

<code>*<-10 ; ->i?</code>	Asynchronous send
---------------------------------	-------------------

GUI application - 1

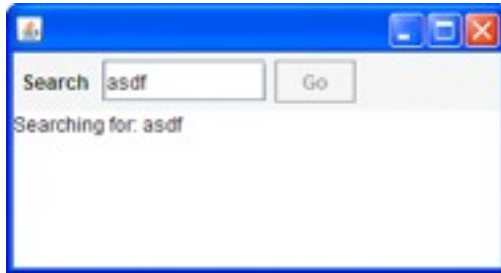


- Input Field
- Search Button
- Searching for...
- Results



GUI application - 2

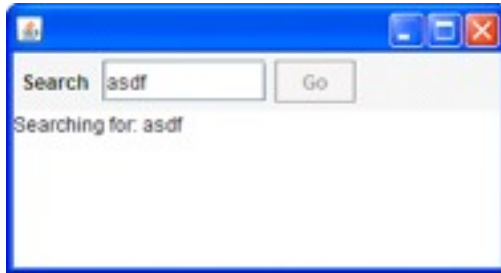
```
val searchButton = new Button("Go") {
  reactions.+= {
    case ButtonClicked(b) =>
      enabled = false
      outputTA.text = "Starting search..."
      new Thread(new Runnable {
        def run() {
          Thread.sleep(3000)
          SwingUtilities.invokeLater(new Runnable{
            def run() {outputTA.text="Search ready"
              enabled = true
            })
          })
        })
      }).start
  }
}
```



GUI application - 3

```
live =      searchButton
           @gui: {outputTA.text="Starting search.."}
                {* Thread.sleep(3000) *}
           @gui: {outputTA.text="Search ready"}
           ...
```

- Sequence operator: white space and ;
- `gui`: code executor for `SwingUtilities.invokeLater+invokeAndWait`
- `{* ... *}`: by executor for `new Thread`



GUI application - 4


live = searchSequence...

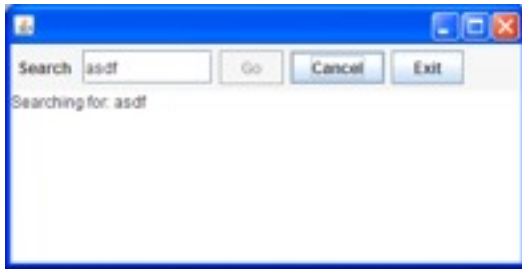
searchSequence = searchCommand
showSearchingText
searchInDatabase
showSearchResults

searchCommand = searchButton
showSearchingText = @gui: {outputTA.text = "..."}
showSearchResults = @gui: {outputTA.text = "..."}
searchInDatabase = {* Thread.sleep(3000) *}

GUI application - 5




- **Search:** button or **Enter** key
- **Cancel:** button or **Escape** key
- **Exit:** button or  **“Are you sure?”**...
- Search only allowed when input field not empty
- Progress indication



GUI application - 6

```

live = searchSequence... || exit

searchCommand = searchButton + Key.Enter
cancelCommand = cancelButton + Key.Escape
exitCommand = exitButton + windowClosing 
exit = exitCommand @gui: while(!areYouSure)
cancelSearch = cancelCommand @gui: showCanceledText

searchSequence = searchGuard searchCommand;
                showSearchingText
                searchInDatabase
                showSearchResults / cancelSearch

searchGuard = if(!searchTF.text.isEmpty) . anyEvent(searchTF) ...

searchInDatabase = {*Thread.sleep(3000)*} || progressMonitor
progressMonitor = {*Thread.sleep( 250)*}
                @gui:{searchTF.text+=here.pass} ...

```


Split Scripts - 1

header: `do~ s:script ~while~ b: =>Boolean ~end`

same as: `do~~while~~end(s:script, b: =>Boolean)`

define: `do~ s:script ~while~ b: =>Boolean ~end = s while(b)`

usage: `test = do~< a;b >~while~ !found ~end`

Split Scripts - 2

```
progressMonitor = sleep_ms(250) updateStatus ...  
                || sleep_ms(5000)
```

```
progressMonitor = during_ms~ 5000  
                  ~every_ms~ 250  
                  ~do~< updateStatus >~end
```

```
during_ms~ duration:Int  
~every_ms~ interval:Int  
~do~ task:script ~end = sleep_ms(interval) task...  
                        || sleep_ms(duration)
```

Script Result Values - 1

```
expr  = term  .. "+"
term  = factor .. "*"
factor = number + "(" expr ")"
```

```
expr  : expr PLUS term { $$ = $1 + $3; }
      | term { $$ = $1; } ;
term  : term MUL factor { $$ = $1 * $3; }
      | factor { $$ = $1; } ;
factor : LPAR expr RPAR { $$ = $2; }
      | NUMBER { $$ = $1; };
```

Script Result Values - 2

```
~ tsk:script ~ f:Unit ~:Int = @onDeactivateWithSuccess{f}: tsk
```

```
expr(?r:Int) = {!r=0!}; var t:Int ~< term(?t)>~~r+=t~ .. "+"
```

```
term(?r:Int) = {!r=1!}; var t:Int ~< factor(?t)>~~r*=t~ .. "*"
```

```
factor(?n:Int) = ?n + "(" expr,?n ")"
```

```
implicit num(??n:Int) = @expNum(_n): {?accept?}
```

Script Result Values - 3

```
~ task: script[Int] ~~ f: Int=>Int ~ : Int
```

```
=
```

```
@onDeactivateWithSuccess{$ = f($task)}: task
```

```
expr : Int = {!0!}^; ~< term >~~ $ + _ ~^ .. "+"
```

```
term : Int = {!1!}^; ~<factor>~~ $ * _ ~^ .. "*"
```

```
factor: Int = ?$ + "(" expr^ ")"
```

Challenges

- **Implementation** 50%: compiler, vm, debugger
- **Unit tests**
- **vms** for simulations, parallel execution, ...
- **New features**
 - **split scripts**
 - **process lambdas**
 - **return values**
 - **data flow**
 - **disambiguation**
- **Documentation, papers, ...**

Challenge: Disambiguation

a b + a c

..a b ; a c

a b || a c

a b |+| a c

..a b |;| a c

Game of Life - 1



Game of Life - 2

```
live           = || boardControl mouseInput speedControl doExit
boardControl   = ...; (...singleStep) multiStep || clear || randomize
doExit        = exitCommand var r=false @gui:{r=areYouSure} while(!r)

randomizeCommand = randomizeButton + 'r'
clearCommand     =   clearButton + 'c'
stepCommand      =   stepButton + ' '
exitCommand      =   exitButton + windowClosing,top
multiStepStartCmd =   startButton + Key.Enter
multiStepStopCmd  =   stopButton + Key.Enter

do1Step        = { *board.calculateGeneration* } @gui: { !board.validate! }

randomize      =   randomizeCommand @gui: { !board.doRandomize()! }
clear          =   clearCommand @gui: { !board.doClear      ! }
singleStep    =   stepCommand do1Step
multiStep     = multiStepStartCmd; ...do1Step { *sleep* }
               / multiStepStopCmd
```

Game of Life - 3

```
speedControl      = ...; speedKeyInput+speedButtonInput+speedSliderInput
```

```
setSpeed(s: Int) = @gui: {!setSpeedValue(s)!}
```

```
speedKeyInput     = times(10)  
                  + val c = chr(pass_up1+'0') key(c)  
                    setSpeed(digit2Speed(c))
```

```
speedButtonInput = if (speed>minSpeed) speedDec  
                  + if (speed<maxSpeed) speedInc
```

```
speedDec          = minSpeedButton setSpeed,minSpeed  
                  + slowerButton  setSpeed(speed-1)
```

```
speedInc         = maxSpeedButton setSpeed,maxSpeed  
                  + fasterButton  setSpeed(speed+1)
```

```
speedSliderInput = speedSlider setSpeed,speedSlider.value
```

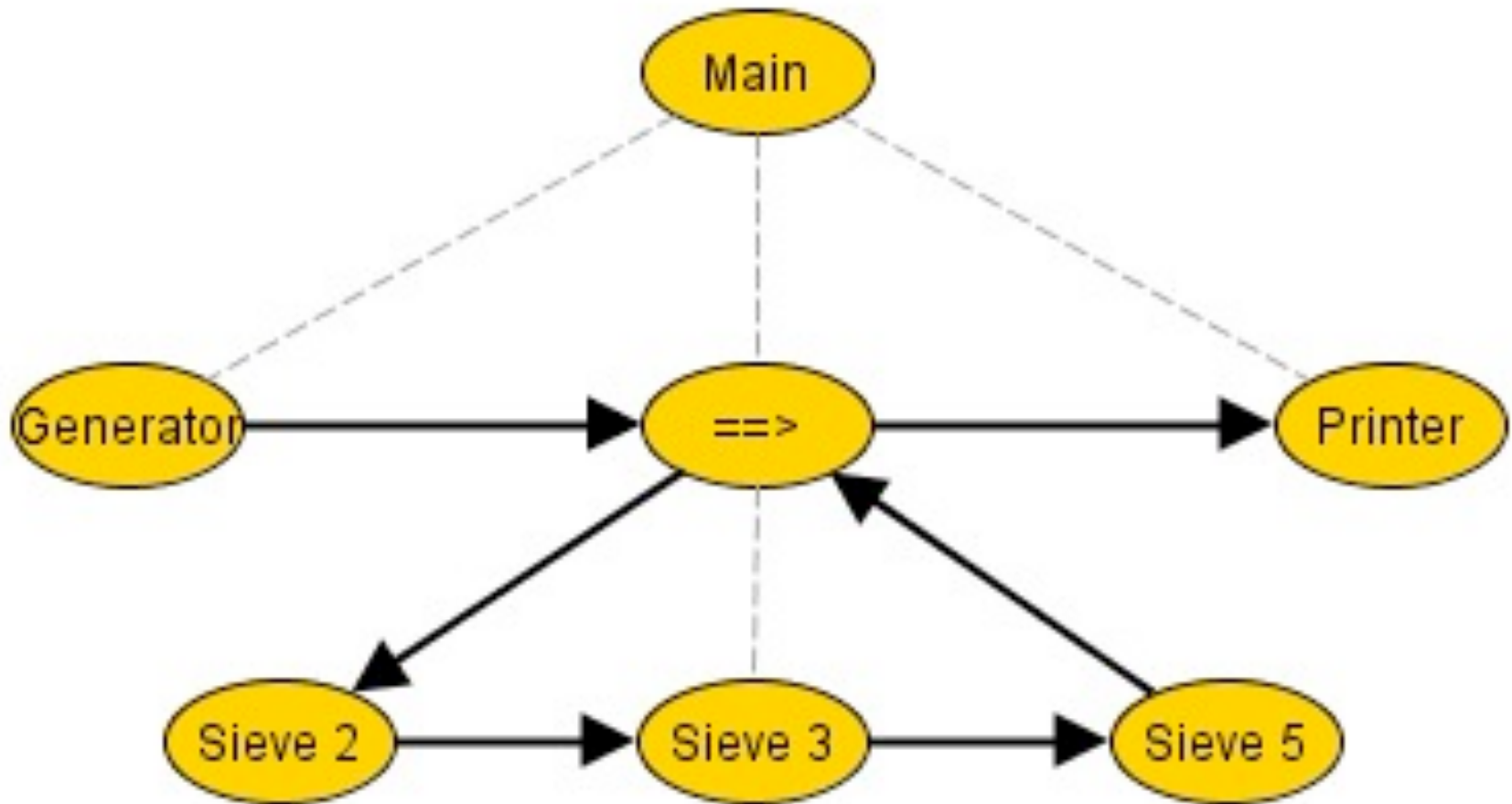
Game of Life - 4

```
mouseInput = (mouseClickInput & mouseDragInput)
/ doubleClick
(mouseMoveInput / doubleClick {!resetLastMousePos!}); ...
```

```
mouseClickInput = var p:java.awt.Point=null
; var doubleClickTimeout=false
mouseSingleClick, board, p?
{! resetLastMousePos !}
( { *sleep_ms(220); doubleClickTimeout=true* }
/ mouseDoubleClick, board, p? )
while (!doubleClickTimeout)
; {! handleMouseSingleClick(p) !}
; ...
```

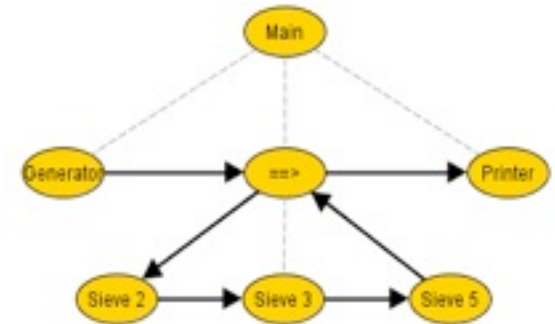
```
mouseMoveInput = mouseMoves( board, (e:MouseEvent)=>handleMove(e.point))
mouseDragInput = mouseDraggings(board, (e:MouseEvent)=>handleDrag(e.point))
/ (mouse_Released {!resetLastMousePos!})
; ...
```

Sieve of Eratosthenes - 1



Sieve of Eratosthenes - 2

```
main = generator(2,1000000)
      ==> (..==>sieve)
      =={toPrint}==> printer
```



```
generator(s:Int,e:Int) = for(i<-s to e) <=i
```

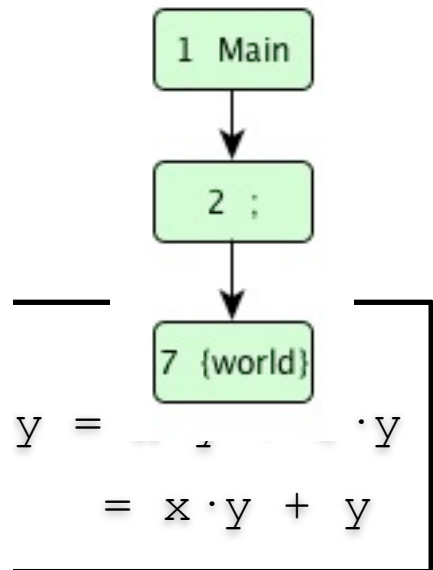
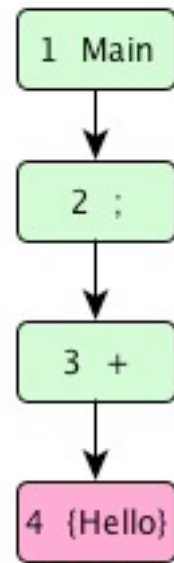
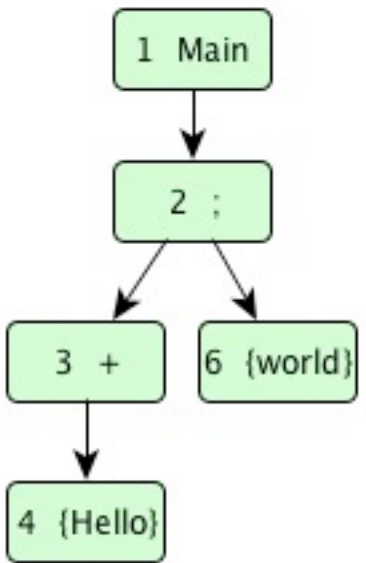
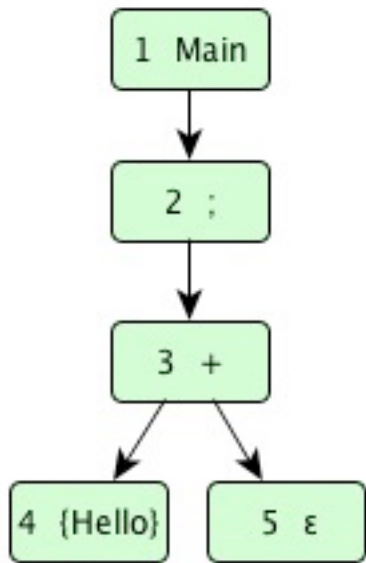
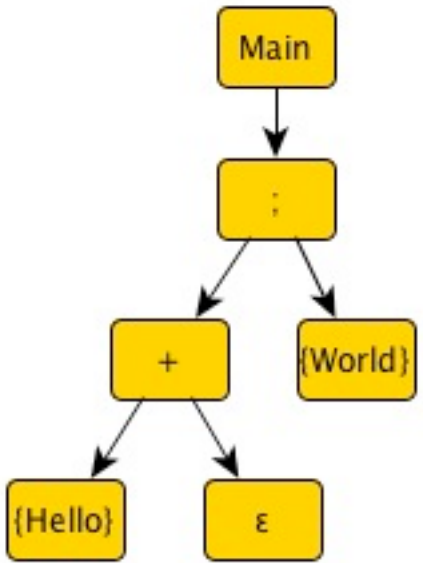
```
sieve      = ==>?p:Int      @toPrint:<=p;
            ..=>?i:Int if (i%p!=0) <=i
```

```
printer    = ..=>?i:Int println,i
```

```
<==>(i:Int) = {}
```

Templates & Call Graphs

{Hello}₊ε; {World}



$$\begin{aligned}
 \bar{y} &= \bar{x} \cdot \bar{y} + \bar{y} \\
 &= x \cdot y + y
 \end{aligned}$$

Experience

- **Scriptic**: Java based predecessor
- In production since 2010
- Analyse technical documentation
- Input: **ODF** ~ **XML** Stream
- **Fun** to use mixture of grammar and 'normal' code
- Parser expectations to scanner

```
implicit text(??s: String) = @expect(here, TextToken(_s): {?accept(here)?})
```

```
implicit number(??n: Int) = @expect(here, NumberToken(_n): {?accept(here)?})
```

- **30,000** accepted of **120,000** expected **tokens per second**

Language Parsing - 1

Low level scripts

`anyText` = `string^`

`anyLine` = `anyText^ endOfLine`

`someEmptyLines` = `..endOfLine`

`someLines` = `{!List[String](C)!}^; .. (anyLine==>{! $:+ _ !}^)`

Language Parsing - 2

For-usage

```
tableRow(ss: String*) = startRow; for(s<-ss) cell(s); endRow
```

```
oneOf(ss: String*) = for(s<-ss) + s^
```

Language Parsing - 3

If-usage

```
footnoteRef(??n: Int) = "(" ??n ")"
```

```
footnote(??n: Int): String = if (fnFormat==NUMBER_DOT) (??n ".")  
                               else (footnoteRef,??n "-")  
; line^  
; .. (line^ ==> {: $ += "."+_.trim :})
```

Experience - 5

Grammar ambiguity

```
var s: String
```

```
var n: Int
```

```
startCell ?s endCell + startCell ?n endCell
```

```
startCell ?s endCell || startCell ?n endCell
```

```
startCell ?s endCell |+| startCell ?n endCell
```

```
xmlTag(t: XMLTag), .. = @expect(here, t) {?accept(here)?}
```